

AN ORE-TYPE CONDITION FOR LARGE k -FACTOR AND DISJOINT PERFECT MATCHINGS

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ABSTRACT. Win [*J. Graph Theory* **6**(1982), 489–492] conjectured that a graph G on n vertices contains k disjoint perfect matchings, if the degree sum of any two nonadjacent vertices is at least $n+k-2$, where n is even and $n \geq k+2$. In this paper, we prove that Win’s conjecture is true for $k \geq n/2$, where n is sufficiently large. To show this result, we prove a theorem on k -factor in a graph under some Ore-type condition. Our main tools include Tutte’s k -factor theorem, the Karush-Kuhn-Tucker theorem on convex optimization, and the solution to the longstanding 1-factor decomposition conjecture.

1. INTRODUCTION

To study the existence of a certain type of subgraphs in a graph is a common topic in graph theory. Maybe the most well-known theorem is the one proved by Dirac [7] in 1952, which is stated as every graph on n vertices has a Hamilton cycle if every vertex of the graph has degree at least $n/2$. Ore [15] extended Dirac’s theorem by considering the degree sum of every pair of nonadjacent vertices in a graph. A graph G is said to be of *Ore-type- (k)* if for every pair of nonadjacent vertices x, y , the degrees of x, y satisfy the inequality $d(x) + d(y) \geq |G| + k$. Ore [16] proved that a graph is *Hamiltonian-connected* if it is of Ore-type-1. Graphs of Ore-type- k were studied by Roberts [17]. Since then, plenty of research was conducted on different graph properties under Ore-type conditions and the variants, such as k -linkedness [10, 13], an equitable coloring of a graph [12], k -ordered Hamiltonicity [9], and etc. Our note mainly concerns on the existence of disjoint perfect matchings in a graph under the Ore-type degree condition.

In 1982, Win [19] posed the following conjecture on disjoint perfect matchings in a graph of Ore-type- $(k-2)$.

Conjecture 1.1 (Win [19]). *Let n, k be two integers such that $1 \leq k \leq n-2$ and n be even. Let G be a simple graph on n vertices. If G is of Ore-type- $(k-2)$, then G contains k disjoint perfect matchings.*

For $k = 1$, Win’s conjecture is true by Ore’s theorem [15]. Win [19] further confirmed the conjecture for $k = 2, 3$.

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On the other hand, the existence of perfect matchings in a graph is closely related to the existence of Hamilton cycles in the same graph. It is an easy observation that every Hamilton cycle in a graph corresponds to a pair of disjoint perfect matchings in the graph, if the order of the graph is even. Egawa [8] proved that: Let $k \geq 2$ be an integer and G be a graph. If $d_G(x) + d_G(y) \geq |G|$ for all non-adjacent vertices x, y , $\delta(G) \geq 2k + 1$, and $|G| \geq 8(2k - 2)^2$, then G has k -edge-disjoint Hamilton cycles. We say that a graph G is a *Fan $2k$ -type graph*, if $d(u, v) = 2$ implies that $\max\{d(u), d(v)\} \geq n/2 + 2k$. Zhou [21] conjectured that every $2k$ -connected Fan $2(k - 1)$ -type graph has k pairwise disjoint Hamilton cycles, and also confirmed this conjecture for $k = 1, 2$. Later, the general case $k \geq 3$ was finally finished by Li [14]. One can easily obtain partial results on Win's conjecture from the results mentioned above.

However, to the best of our knowledge, Win's conjecture is still wide open now. One of our results concerns Win's conjecture when k is large in compare with n .

Theorem 1.2. *Win's conjecture is true for sufficiently large even n , if $k \geq n/2$.*

In this paper, instead of proving Theorem 1.2 directly, we firstly prove our main result which focuses on the existence of large k -factors.

Theorem 1.3. *Let n and k be two integers such that $n \geq k + 1 \geq n/2 + 1$ and kn be even. Let G be a graph on n vertices. If G is of Ore-type- $(k - 2)$, then G contains a k -factor.*

With the help of Theorem 1.3, we will use the solution to 1-factor decomposition conjecture to prove Theorem 1.2. Recall that the long-standing 1-factorization conjecture states that every regular graph of sufficiently large degree has a 1-factorization. It was first stated explicitly by Chetwynd and Hilton [4, 5], and they also stated by Dirac, who discussed it in the 1950s. Partial results were obtained by Chetwynd and Hilton [4, 5], and Zhang and Zhu [20]. Recently, Csaba et al. [6] confirmed this conjecture for large graphs. One of their main results in [6] is used for our proof of Theorem 1.2.

Theorem 1.4 (Csaba et al. [6]). *Suppose that n is sufficiently large and even, and $D \geq 2\lceil n/4 \rceil - 1$. Then every D -regular graph G on n vertices has a decomposition into perfect matchings.*

The proof of our main theorem also uses a theorem of Katerinis and Woodall on k -factor, and the Karush-Kuhn-Tucker theorem on convex optimization. We will introduce all necessary terminology and additional results in the next section.

Now we give some necessary notation and terminology. Let G be a graph. We use $V(G)$ and $E(G)$ to denote the vertex set and edge set of G , respectively, and denote by $|G| = |V(G)|$. Let S, T be two disjoint subsets of $V(G)$, $E_G(S, T)$ be the set of edges between S and T in G , and $e_G(S, T) = |E_G(S, T)|$. When S consists of a single element, say $S = \{v\}$, we use $E_G(v, T)$ and $e_G(v, T)$ instead of $E_G(\{v\}, T)$ and $e_G(\{v\}, T)$, respectively. Let $v \in V(G)$ and H be a subgraph of G . $N_G(v)$ is the set of neighbors of v in G and $d_G(v) = |N_G(v)|$. Set $N_H(v) = N_G(v) \cap V(H)$ and $d_H(v) = |N_H(v)|$. When

there is no danger of ambiguity, we use $d(v)$ instead of $d_G(v)$ for short. Let $S \subset V(G)$ and let $G - S$ denote the subgraph of G induced by $V(G) \setminus V(S)$. If S consists of only one vertex, say $S = \{v\}$, we use $G - v$ instead of $G - \{v\}$. For notation and terminology not defined here, we refer the reader to Bondy and Murty [2].

The organization of our paper is as follows. In Section 2, we introduce necessary preliminaries. In Section 3, we prove Theorems 1.2 and 1.3.

2. SOME PRELIMINARIES

In this section, we first introduce some notation and terminology related to Tutte's k -factor theorem. For any pair of disjoint subsets $S, T \subset V(G)$, a component C of $G - S - T$ is called a k -odd-component if

$$e_G(V(C), T) + k|V(C)| \equiv 1 \pmod{2}.$$

We usually use $q(S, T)$ to denote the number of components of $G - S - T$ which are k -odd components.

Tutte's k -factor theorem is well known.

Theorem 2.1 (Tutte [18]). *Let k be a positive integer. A graph G contains no k -factor if and only if there exist disjoint subsets $S, T \subset V(G)$, such that*

$$(2.1) \quad \eta(S, T) := k|S| - k|T| + \sum_{x \in T} d_{G-S}(x) - q(S, T) \leq -2.$$

From Tutte's theorem, Katerinis and Woodall [11] deduced the following. It shall play an important role in our proof.

Theorem 2.2 (Katerinis and Woodall [11]). *Let $k \geq 1$ be an integer. If a graph G contains no k -factor, then there exist two disjoint subsets $S, T \subset V(G)$ such that there holds (2.1), and*

$$(2.2) \quad e_G(v, T) \leq k - 1, \text{ and}$$

$$(2.3) \quad d_{G-S}(v) \geq k + 1 \text{ for all } v \in U,$$

where U denotes the union of all k -odd components of $G - S - T$.

Our proof also uses tools from optimization. An optimization problem of the form

$$(2.4) \quad \begin{cases} \min & f(x), \\ \text{s.t.} & g_i(x) \leq 0, \text{ for } i = 1, \dots, m \end{cases}$$

is called a *convex optimization problem* if the functions $f, g_1, \dots, g_m : R_n \rightarrow R$ are all convex functions. We need the Karush-Kuhn-Tucker theorem on convex optimization. The following one is a direct corollary of Theorem 4.3.8 in [1, pp.207].

Theorem 2.3 (Karush-Kuhn-Tucker sufficient condition [1]). *Let X be a nonempty open set in R^n , and let $f : R^n \rightarrow R$, $g_i : R^n \rightarrow R$ for $i = 1, \dots, m$. Consider Problem P :*

$$(2.5) \quad \begin{cases} \min & f(x), \\ \text{s.t.} & g_i(x) \leq 0, \text{ for } i = 1, \dots, m \\ & x \in X, \end{cases}$$

Let \bar{x} be a local optimal solution. There exist scalars $u_i \geq 0$ for $1 \leq i \leq m$ such that

$$(2.6) \quad \nabla f(\bar{x}) + \sum_{i \in I} u_i \nabla g_i(\bar{x}) = 0.$$

The point satisfying (2.6) is called a *KKT point*. For convex optimal problems, the KKT conditions are also sufficient for optimality (see [1, pp.773].).

Theorem 2.4. *For the convex optimal problem (2.4), every KKT point is a global optimal solution.*

The next result is a well-known result on convex function.

Theorem 2.5. *Let $f(x)$ be a function on R , where R is a convex set. Suppose that f is twice differentiable and f'' is continuous. Then $f(x)$ is a convex function if and only if its Hessian matrix is positive semi-definite on R .*

For more information and details, we refer the reader to Boyd and Vandenberghe [3].

3. PROOFS OF THEOREMS 1.2 AND 1.3.

In this section, we will present the proofs of Theorems 1.2 and 1.3.

Proof of Theorem 1.3. We prove Theorem 1.3 by contradiction. Suppose that G contains no k -factors. By Theorem 2.2, we can choose disjoint $S, T \subset V(G)$ satisfying (2.1), (2.2), and (2.3). Define $s := |S|$ and $t := |T|$. Let C_1, \dots, C_q be all k -odd components of $G - S - T$. So, for every vertex $v \in V(C_i)$, $d_{G-S}(v) \geq k + 1$ and $e_G(v, T) \leq k - 1$, and this implies $d_{C_i}(v) \geq 2$. Thus, $|C_i| \geq 3$.

Claim 1. *G is k -connected, and hence, the minimum degree $\delta(G) \geq k$.*

Proof. Let W be a cut-set of G and let C'_1, C'_2 be two components of $G - W$. For $x \in V(C'_1)$ and $y \in V(C'_2)$, one can see that $xy \notin E(G)$, and thus

$$n + k - 2 \leq d(x) + d(y) \leq |C'_1| + |C'_2| - 2 + 2|W|.$$

Notice that $n \geq |C'_1| + |C'_2| + |W|$. Hence, $|W| \geq k$, and moreover, $\delta(G) \geq k$. \square

Now we show that $T \neq \emptyset$. Otherwise, by (2.1) and Claim 1, we have $q(S, \emptyset) \geq ks + 2 \geq k^2 + 2$. Thus, $n \geq |U| + s + t \geq 3(k^2 + 2) + k \geq \frac{3}{4}n^2 + \frac{1}{2}n + 6$, which is impossible.

Set $h_1 := \min\{d_{G-S}(x) : x \in T\}$. Let $u_1 \in T$ such that $d_{G-S}(u_1) = h_1$. Set $N_T[u_1] := (N(u_1) \cap T) \cup \{u_1\}$. For any vertex $x \in V(G)$, let $d_T(x) = |N_G(x) \cap T|$. If $T - N_T[u_1] \neq \emptyset$, let $h_2 := \min\{d_{G-S}(x) : x \in T - N_T[u_1]\}$ and choose $u_2 \in T - N_T[u_1]$ such that $d_{G-S}(u_2) = h_2$. As in the proof of Lemma 2.2, we still denote $U := C_1 \cup C_2 \dots \cup C_q$.

Claim 2.

$$(3.1) \quad s + h_1 \geq k.$$

Proof. Since $\delta(G) \geq k$, $s + h_1 \geq d_G(u_1) \geq k$. \square

In the following, we divide the proof into four cases.

Case 1. $h_1 \geq k$.

By (2.1), we have

$$q := q(S, T) \geq k|S| - k|T| + 2 + \sum_{x \in T} d_{G-S}(x) \geq k|S| - k|T| + 2 + h_1|T| \geq ks + 2 \geq 2.$$

This means that $G - S - T$ is disconnected. By Claim 1, $s + t \geq k$. Notice that $k \geq n/2$. Since $|C_i| \geq 3$ for each $i = 1, \dots, q$, we infer that $|U| \geq 3q \geq 3(ks + 2)$. If $|S| \geq 1$, then $n = |G| \geq |U| + s + t \geq 3(k + 2) + s + t \geq 4k + 6 > n$, a contradiction. Thus, $S = \emptyset$ and $t \geq k$. Since $q \geq 2$, choose $x \in V(C_1)$ and $y \in V(C_2)$, and we have

$$\begin{aligned} n + k - 2 &\leq d(x) + d(y) \\ &\leq |C_1| - 1 + |N_G(x) \cap T| + |C_2| - 1 + |N_G(y) \cap T| \\ &\leq |C_1| + |C_2| + 2k - 4 \quad (\text{by Theorem 2.2}) \\ &\leq n - t + 2k - 4 \\ &\leq n + k - 4, \end{aligned}$$

a contradiction.

Thus, in the following, assume that

$$(3.2) \quad h_1 \leq k - 1.$$

Case 2. $T = N_T[u_1]$.

Claim 3. For any $i \in \{1, \dots, q\}$, there exists $w_i \in V(C_i)$ such that $w_i u_1 \notin E(G)$

Proof. Suppose that there exists $j \in \{1, \dots, q\}$, such that $V(C_j) \subset N_{G-S}(u_1)$. Notice that for $x \in V(C_j)$, $d_{G-S}(x) \geq k + 1$, and $N_{G-S}(x) \subset V(C_j) \cup T$. Then by (3.2), $k - 1 \geq h_1 = d_{G-S}(u_1) \geq |C_j| + |T| - 1 \geq d_{G-S}(x) \geq k + 1$, a contradiction. \square

Claim 4.

$$(3.3) \quad |C_i| \geq k - h_1 + 2.$$

$$(3.4) \quad n \geq s + t + q(k - h_1 + 2).$$

Proof. For each $i \in \{1, 2, \dots, q\}$, by Claim 3, there exists a vertex $x_i \in V(C_i)$ such that $x_i u_1 \notin E(G)$. Since $d_T(x_i) \leq |T| - 1 = d_T(u_1) \leq d_{G-S}(u_1) = h_1$, we have

$$\begin{aligned} |C_i| &\geq |N_G(x_i) \cap V(C_i)| + 1 \\ &= d_{G-S}(x_i) - d_T(x_i) + 1 \end{aligned}$$

$$\begin{aligned}
&\geq (k+1) - h_1 + 1 \\
&= k - h_1 + 2.
\end{aligned}$$

Moreover, by (3.3), we can get

$$n = |G| \geq |S| + |T| + \sum_{i=1}^q |C_i| \geq s + t + q(k - h_1 + 2).$$

□

Claim 5. $q = q(S, T) \geq 2$.

Proof. By (3.2), the fact $T = N_T[u_1]$, and the definition of u_1 , we infer

$$d_T(u_1) = t - 1 \leq d_{G-S}(u_1) = h_1 \leq k - 1,$$

Thus, $k \geq h_1 + 1 \geq t$. So,

$$\begin{aligned}
(3.5) \quad q(S, T) &\geq ks - kt + \sum_{x \in T} d_{G-S}(x) + 2 \\
&\geq ks - kt + h_1 t + 2
\end{aligned}$$

$$\begin{aligned}
&\geq k(k - h_1) - kt + h_1 t + 2 \\
&= (k - h_1)(k - t) + 2 \\
(3.6) \quad &\geq 2.
\end{aligned}$$

□

Claim 6.

$$(3.7) \quad s \geq k + (q - 1)(k + 2 - h_1) + t - 2h_1 - 1.$$

Proof. For any $i \in \{1, \dots, q\}$, since $w_i u_1 \notin E(G)$, we have

$$(3.8) \quad d(w_i) + d(u_1) \geq n + k - 2.$$

On the other hand, we obtain

$$d_T(w_i) \leq t - 1 = d_T(u_1) \leq d_{G-S}(u_1) = h_1.$$

One can see that

$$(3.9) \quad d(w_i) + d(u_1) \leq |C_i| - 1 + 2h_1 + 2s.$$

Combining (3.8) and (3.9), we can infer

$$\begin{aligned}
&n + q(2h_1 + 2s - 1) \\
&\geq \sum_{i=1}^q |C_i| + q(2h_1 + 2s - 1) + s + t \\
&\geq q(n + k - 2) + s + t,
\end{aligned}$$

that is,

$$(q-1)n \leq q(2h_1 + 2s - 1 - k + 2) - s - t.$$

By Claim 5, $q \geq 2$. By (3.4), we have

$$(q-1)(s+t+q(k-h_1+2)) \leq q(2h_1+2s-k+1) - s - t.$$

This implies

$$(q-1)(k+2-h_1) + s+t \leq 2h_1+2s-k+1,$$

and this proves the claim. \square

By computation, we have

$$\begin{aligned} 0 &\geq 2 + ks - kt + \sum_{x \in T} d_{G-S}(x) - q \\ &\geq 2 + k(k + (q-1)(k+2-h_1) - 2h_1 - 1) + h_1t - q \quad (\text{by (3.7)}) \\ &= 2 + q(k(k+2-h_1) - 1) - k(3+h_1) + h_1t \\ &\geq 2 + (ks - kt + h_1t + 2)(k(k+2-h_1) - 1) - k(3+h_1) + h_1t \quad (\text{by (3.2) and (3.5)}) \\ &\geq 2 + (ks - (k-h_1)(h_1+1) + 2)(k(k+2-h_1) - 1) \\ &\quad - k(3+h_1) + h_1(h_1+1) \\ &\geq 2 + 2(k(k+2-h_1) - 1) - k(3+h_1) + h_1(h_1+1) \\ &= 2k^2 + k - 3kh_1 + h_1^2 + h_1 \\ &\geq 3k, \end{aligned}$$

where we have used the fact $k(k+2-h_1) - 1 \geq 3k - 1 \geq 0$ in the third inequality above; and (3.1), (3.2) and (3.5) in the fifth inequality above; and the fact $f(h_1) \geq f(k-1)$ in the last step, where the function $f(x) = -3kx + x^2 + x$, $x \leq k-1$.

This contradiction completes the proof of the case.

Case 3. $T \neq N_T[u_1]$ and $h_2 \geq k$.

Set $p := |N_T[u_1]|$. Recall that $V(U) = V(C_1 \cup \dots \cup C_q)$. We have

Claim 7. $q(S, T) \geq 2$, where the equality holds when $h_1 = k-1$, $p = k$ and $h_2 = k$.

Proof. By (2.1), we have

$$\begin{aligned} q(S, T) &\geq k|S| - k|T| + \sum_{x \in T} d_{G-S}(x) + 2 \\ &\geq ks - kt + h_1p + h_2(t-p) + 2. \end{aligned}$$

By the hypothesis $h_2 \geq k$ and $t \geq p$, we obtain

$$(3.10) \quad ks - kt + h_1p + h_2(t-p) + 2 \geq ks - (k-h_1)p + 2.$$

By (3.1), $s \geq k - h_1$. Since $s \geq k - h_1$ and $p \leq h_1 + 1$, we obtain $ks - (k-h_1)p + 2 \geq (k-h_1)(k-h_1-1) + 2$. By (3.2), $k \geq h_1 + 1$. So, $(k-h_1)(k-h_1-1) + 2 \geq 2$, and

hence $q(S, T) \geq 2$. The condition when the equality holds can be deduced easily. This proves the claim. \square

Claim 8. Suppose there exists a vertex $x \in V(U)$ such that $xu_1 \notin E(G)$. Then

$$(3.11) \quad s \geq k + 3(q - 1) - h_1.$$

Proof. Without loss of generality, assume that $x \in V(C_i)$ for some $i \in \{1, \dots, q\}$. By Lemma 2.2, one may see that $|C_i| \geq 3$ for $i = 1, \dots, q$. So we obtain

$$n \geq s + t + 3q.$$

We also have

$$n + k - 2 \leq d(x) + d(u_1) \leq (|C_i| - 1) + (t - 1) + s + h_1 + s = 2s + t + |C_i| + h_1 - 2.$$

One can see that

$$|C_i| \leq n - s - t - 3(q - 1).$$

Thus, we have

$$n + k - 2 \leq s + n + h_1 - 3(q - 1) - 2.$$

This proves the claim. \square

Claim 9. $V(U) \subset N_G(u_1)$.

Proof. Suppose not. By Claim 8, (3.11) holds. Thus,

$$\begin{aligned} 0 &\geq 2 + ks - kt + \sum_{x \in T} d_{G-S}(x) - q \\ &\geq 2 + k(3(q - 1) + k - h_1) - kt + h_1p + h_2(t - p) - q \quad (\text{by (3.11)}) \\ &\geq 2 + k(3(q - 1) + k - h_1) + (h_1 - k)(h_1 + 1) - q \\ &\geq k(3 + k - h_1) + (h_1 - k)(h_1 + 1) \\ &= h_1^2 - (2k - 1)h_1 + k^2 + 2k \\ &\geq 3k, \end{aligned}$$

a contradiction. Notice that in the above, we have used the facts $h_2 \geq k$, $t \geq p$, $h_1 \leq k - 1$ and $p \leq h_1 + 1$ in the third step; and the facts that the function $f(q) = 3k(q - 1) - q$ is increasing and $q \geq 2$ (by Claim 7) in the fourth step; and the fact that the function $f(h_1) = h_1^2 - (2k - 1)h_1 + k^2 + 2k$ is decreasing when $h_1 \leq k - 1$ in the last step.

The proof of this claim is complete. \square

By Claim 9, $V(U) \subset N_G(u_1)$. So, $h_1 \geq 3q + p - 1$. We have

$$\begin{aligned} 0 &\geq 2 + ks - kt + \sum_{x \in T} d_{G-S}(x) - q \\ &\geq 2 + k(k - h_1) - kt + h_1p + h_2(t - p) - q \quad (\text{by the fact } s + h_1 \geq k) \\ &\geq 2 + k(k - h_1) + (h_1 - k)p - q \quad (\text{by (3.10)}) \end{aligned}$$

$$\begin{aligned}
&\geq 2 + (k - h_1)(k - p) - q \\
&\geq 2 + (k - h_1)(k - h_1 + 3q - 1) - q \\
&\geq (k - h_1)(k - h_1 + 5) \quad (\text{since } q \geq 2) \\
&> 0,
\end{aligned}$$

a contradiction. This proves the case.

Case 4. $0 \leq h_1 \leq h_2 \leq k - 1$.

Since $u_1 u_2 \notin E(G)$, it follows that

$$n + k - 2 \leq d(u_1) + d(u_2) \leq h_1 + h_2 + 2s,$$

i.e.,

$$(3.12) \quad s \geq \frac{1}{2}(n + k - 2 - h_1 - h_2).$$

Since $|C_i| \geq 3$, one may see that

$$(3.13) \quad n \geq s + t + 3q.$$

We can get

$$\begin{aligned}
0 &\geq ks - kt + h_1 p + h_2(t - p) + 2 - q \\
&= ks - (k - h_2)t + (h_1 - h_2)p + 2 - q \\
&\geq ks - (k - h_2)(n - s - 3q) + (h_1 - h_2)p + 2 - q \quad (\text{by (3.13)}) \\
&\geq (2k - h_2)s - (k - h_2)n + q(3(k - h_2) - 1) + (h_1 - h_2)(h_1 + 1) + 2 \quad (\text{since } p \leq h_1 + 1, h_1 \leq h_2) \\
&\geq (2k - h_2)s - (k - h_2)n + (h_1 - h_2)(h_1 + 1) + 2,
\end{aligned}$$

i.e.,

$$(3.14) \quad 0 \geq (2k - h_2)s - (k - h_2)n + (h_1 - h_2)(h_1 + 1) + 2.$$

First suppose that

$$(3.15) \quad h_1 - h_2 \geq k + 2 - n.$$

One can see that

$$\begin{aligned}
0 &\geq \frac{1}{2}(2k - h_2)(n + k - 2 - h_1 - h_2) - (k - h_2)n + (h_1 - h_2)(h_1 + 1) + 2 \quad (\text{by (3.12)}) \\
&= h_1^2 - h_1(k - 1 + \frac{h_2}{2}) + \frac{h_2^2}{2} + \frac{1}{2}(n - 3k)h_2 + k^2 - 2k + 2 \\
&\geq h_1^2 - h_1(k - 1 + \frac{h_2}{2}) + \frac{h_2^2}{2} + \frac{1}{2}((k + 2 + h_2 - h_1) - 3k)h_2 + k^2 - 2k + 2 \quad (\text{by (3.15)}) \\
&= h_1^2 - h_1(k - 1 + h_2) + h_2^2 + (-k + 1)h_2 + k^2 - 2k + 2,
\end{aligned}$$

i.e.,

$$(3.16) \quad 0 \geq h_1^2 - h_1(k - 1 + h_2) + h_2^2 + (-k + 1)h_2 + k^2 - 2k + 2.$$

Let $f(h_1, h_2, k) = h_1^2 - h_1(k - 1 + h_2) + h_2^2 + (-k + 1)h_2 + k^2 - 2k + 2$. Consider the following non-linear programming problem:

$$(3.17) \quad \begin{cases} \min & f(h_1, h_2, k), \\ \text{s.t.} & h_1 - h_2 \leq 0, \\ & h_2 \leq k - 1, \\ & -h_1 \leq 0, \end{cases}$$

The Hessian matrix of the function $f(h_1, h_2, k)$ is

$$M = \begin{pmatrix} 2 & -1 & -1 \\ -1 & 2 & -1 \\ -1 & -1 & 2 \end{pmatrix}.$$

Note that M is a positive semi-definite matrix. By Theorem 2.5, $f(h_1, h_2, k)$ is a convex function. Thus (3.17) is a convex optimization problem. Its Lagrangian function is

$$\begin{aligned} L(h, \lambda) = & h_1^2 - h_1(k - 1 + h_2) + h_2^2 + (-k + 1)h_2 + k^2 - 2k + 2 + \lambda_1(h_1 - h_2) \\ & + \lambda_2(h_2 - k + 1) + \lambda_3(-h_1). \end{aligned}$$

Hence the Karush-Kuhn-Tucker condition of (3.17) is

$$(3.18) \quad \begin{cases} 2h_1 - (k - 1 + h_2) + \lambda_1 - \lambda_3 = 0, \\ -h_1 + 2h_2 + (-k + 1) - \lambda_1 + \lambda_2 = 0, \\ -h_1 - h_2 + 2k - 2 - \lambda_2 = 0, \\ \lambda_1(h_1 - h_2) = 0, \\ \lambda_2(h_2 - k + 1) = 0, \\ \lambda_3 h_1 = 0. \end{cases}$$

It is easy to see that $h_1 = h_2 = k - 1$ and $\lambda_1 = \lambda_2 = \lambda_3 = 0$ is a solution of the equation (3.18). For a convex optimization problem, by Theorem 2.3, every solution satisfying its KKT condition is also its optimum solution. Thus, we have

$$f(h_1, h_2, k) \geq f(k - 1, k - 1, k) = 1,$$

contradicting (3.16).

Finally, suppose that

$$h_2 - h_1 > n - k - 2.$$

By (3.14) and (3.1), one can see that

$$\begin{aligned} 0 & \geq (2k - h_2)s - (k - h_2)n + (h_1 - h_2)(h_1 + 1) + 2 \\ & \geq (2k - h_2)(k - h_1) - (k - h_2)n + (h_1 - h_2)(h_1 + 1) + 2 \\ & = h_1^2 - (2k - 1)h_1 + 2k^2 - kn + 2 + h_2(n - k - 1) \\ & \geq h_1^2 - (2k - 1)h_1 + 2k^2 - kn + 2 + (n - k - 1)^2 + (n - k - 1)h_1 \\ & = h_1^2 - (3k - n)h_1 + 2k^2 - kn + 2 + (n - k - 1)^2 \\ & \geq -\frac{1}{4}(3k - n)^2 + 2k^2 - kn + 2 + (n - k - 1)^2 \end{aligned}$$

$$\begin{aligned}
&= -\frac{1}{4}(n-k)^2 + (n-k-1)^2 + 2 \\
&= \frac{3}{4}(n-k)^2 - 2(n-k) + 3 \\
&> 0,
\end{aligned}$$

a contradiction. This completes the proof of Theorem 1.3. \square

Proof of Theorem 1.2. By Theorem 1.3, G contains a k -factor, denoted by H , where $k \geq n/2 \geq 2\lceil n/4 \rceil - 1$. Obviously, H is k -regular. Since the order of G is sufficiently large, the order of H is also sufficiently large. By Theorem 1.4, H can be decomposed into k disjoint perfect matchings. The proof of Theorem 1.2 is completed. \square

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